

Entity Authentication

Prof. Bart Preneel
 COSIC – KU Leuven - Belgium
 Firstname.Lastname(at)esat.kuleuven.be
 http://homes.esat.kuleuven.be/~preneel
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Goals

- Understand goals of entity authentication
- Understand strength and limitations of entity authentication protocols including passwords
- Understand subtle problems when entity authentication protocols are deployed in practice

Definitions (ctd)

		data	
Confidentiality	confidentiality	encryption	anonymity
Integrity			
Availability	authentication	data authentication	identification

Authorisation

Non-repudiation of origin, receipt

Contract signing

Notarisation and Timestamping

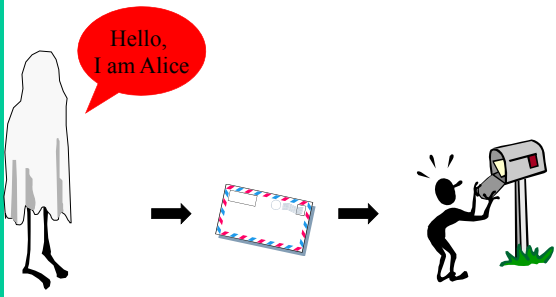
E-voting, e-auction,...

Don't use the word authentication without defining it

Identification

- the problem
- passwords
- challenge response with symmetric key and MAC (symmetric tokens)
- challenge response with public key (signatures, ZK)
- biometry

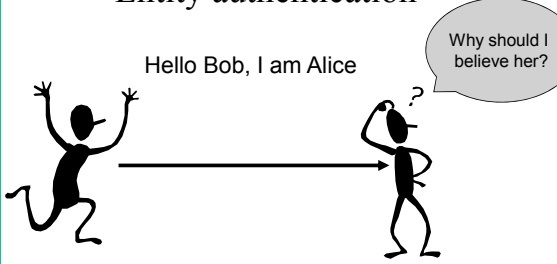
Entity authentication



Hello, I am Alice

Eve
Bob

Entity authentication



Hello Bob, I am Alice

Why should I believe her?

entity authentication: one is corroborated of the identity of another party, and of the fact that this party is **alive (active)** during the protocol

Entity authentication is based on one or more of the following elements:

- what someone **knows**
 - password, PIN
- what someone **has**
 - magstripe card, smart card
- what someone **is** (biometrics)
 - fingerprint, retina, hand shape,...
- **how** someone does something
 - manual signature, typing pattern
- **where** someone is
 - dialback, location based services (GSM, Galileo)

ert5^r\$#890y



Entity authentication with passwords

BUT

- Eve can guess the password
- Eve can listen to the channel and learn Alice's password
- Bob needs to know Alice's secret
- Bob needs to store Alice's secret in a secure way

Possibility of replay: liveliness is missing

Improved identification with passwords

Bob stores $f(P)$ rather than Alice's secret P

- it is difficult to deduce P from $f(P)$

Password entropy: effective key length

Category	5 chars	6 chars	7 chars	8 chars	9 chars	10 chars
lower case	~25	~28	~32	~35	~38	~42
lower case + digits	~28	~32	~35	~38	~42	~45
mixed case+digits	~32	~35	~38	~42	~45	~48
keyboard	~35	~38	~42	~45	~48	~52

Problem: passwords from dictionaries

Improved+ identification with passwords

Bob stores $f(P,S) || S$ rather than Alice's secret P

it is harder to attack the passwords of all users simultaneously

give every user at registration a random publicly known value S (salt)

Example: UNIX


- Function $f()$ = DES applied 25 times to the all zero plaintext with as key the password P (8 7-bit characters)
- Salt: 12-bit modification to DES
- etc/passwd public
- PC: 20-40 million passwords/second
- But time-memory tradeoff...
 - Precomputation per salt $25 \cdot 2^{56}$
 - Storage per salt: 2 Terabyte
 - Find one key in time $25 \cdot 2^{38}$

Improving password security

- Apply the function f “ x ” times to the password (iteratively)
 - if $x = 100$ million, testing a password guess takes a few seconds
 - need to increase x with time (Moore’s law)
 - Examples: PBKDF2 (Password-Based Key Derivation Function 2), scrypt, bcrypt
- Disadvantage: one cannot use the same hashed password file on a faster server and on an embedded device with an 8-bit microprocessor
 - need to use different values of x depending on the computational power of the machine

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Problem: human memory is limited



- Solution: store key K on magstripe, USB key, hard disk
- Stops guessing attacks

But this does not solve the other problems related to passwords
And now you identify the card, not the user....


Possibility of replay: liveness is missing

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Improvement: Static Data Authentication

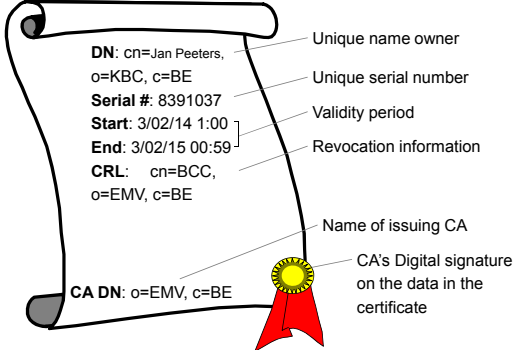
- Replace K by a signature of a third party CA (Certification Authority) on Alice’s name: $\text{SigSK}_{CA}(\text{Alice})$ = special certificate
- Advantage: can be verified using a public string PK_{CA}
- Advantage: can only be generated by CA
- Disadvantage: signature = 40..128 bytes
- Disadvantage: can still be copied/intercepted

Possibility of replay: liveness is missing



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“Certificate” for static data authentication

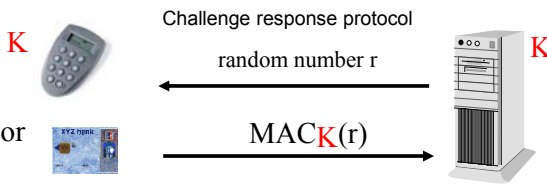


DN: cn=Jan Peeters, o=KBC, c=BE
Serial #: 8391037
Start: 3/02/14 1:00
End: 3/02/15 00:59
CRL: cn=BCC, o=EMV, c=BE
CA DN: o=EMV, c=BE

- Unique name owner
- Unique serial number
- Validity period
- Revocation information
- Name of issuing CA
- CA’s Digital signature on the data in the certificate

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Entity authentication with symmetric token



Challenge response protocol

random number r

$\text{MAC}_K(r)$

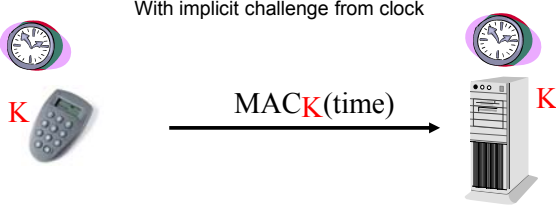
- Eavesdropping no longer effective
- Bob still needs secret key K

Detects whether Alice is alive!

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Entity authentication with symmetric token

With implicit challenge from clock



$\text{MAC}_K(\text{time})$

- Eavesdropping no longer effective
- Bob still needs secret key K
- resynchronization mechanism needed

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Lamport's one-time passwords

iterated one-way function

• Disadvantage: only works with one Bob

Entity authentication with public key token

Challenge response protocol

- Eavesdropping no longer effective
- Bob no longer needs a secret – only PK_A

Entity authentication with ZK

Zero knowledge

- Mathematical proof that Bob only learns that he is talking to Alice (1 bit of information)
- Bob cannot use this information to convince a third party that he is/was talking to Alice

ZK definitions

- **complete:** if Alice knows the secret, she can carry out the protocol successfully
- **sound:** Eve (who wants to impersonate Alice) can only convince Bob with a very small probability that she is Alice;
- **zero knowledge:** even a dishonest Bob does not learn anything except for 1 bit (he is talking to Alice); he could have produced himself all the other information he obtains during the protocol.

ZK: Fiat-Shamir (1986)

- central RSA modulus n
- per user:
 - identity I_A
 - secret key s_A ($0 < s_A < n$)
 - public key $y_A = s_A^2 \pmod n$
- facts from number theory:
 - if one knows the factorization of n , it is easy to compute the square roots modulo n (if they exist);
 - if one can compute square roots modulo n , it is easy to factor n

ZK: Fiat-Shamir

All operations mod n

• **Complete:** trivial

• **Sound:** Eve's probability of success = $\frac{1}{2}$
 Eve gambles that Bob will choose $e=0$
 then she chooses r , and computes $x=r^2$ and $z_0=r$
 Eve gambles that Bob will choose $e=1$
 then she chooses z_1 , and computes $x=z_1^2/y_A$
 If Eve knows both z_0 and z_1 then she knows $s_A = z_1/z_0$

ZK: Fiat-Shamir

- **zero knowledge:** Bob learns nothing about Alice's secret
- $e=0$: B sees r and r^2
- $e=1$: B sees r^2 (from $r^2 s_A^2 = r^2 \cdot y_A$) and $r s_A$
 - $r \cdot s_A$ is a Vernam encryption of s_A : statistically independent of s_A
- Hence B only sees 2 random squares mod n , which he could have produced himself (yet he is convinced that he has spoken to Alice!)
- in practice: more iterations (20...40) for better security ($1/2^{20} \dots 1/2^{40}$)

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Overview Identification Protocols

	Guess	Eavesdrop channel (liveliness)	Impersonation by Bob	Secret info for Bob	Security
Password	-	-	-	-	1
Magstripe (SK)	+	-	-	-	2
Magstripe (PK)	+	-	-	+	3
Dynamic password	+	+	-	-	4
Smart card (SK)	+	+	-	-	4
Smart Card (PK)	+	+	+	+	5

Entity authentication with password

Challenge response protocol

- Eavesdropping no longer effective
- Bob still needs secret key P
- Exhaustive search for P is easy based on a single transcript

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Entity authentication with password: EKE

[Bellovin, Merritt '92]
All operations mod p

$x \in_R [1, p-1]$
 r_A 128-bit string
 $k = (\alpha^y)^x$

$y \in_R [1, p-1]$
 r_B 128-bit string
 $k = (\alpha^x)^y$

- Adds entity authentication to Diffie Hellman
- Attacker cannot perform off-line exhaustive search for the password P
- Attacker can still try on-line attacks; need to restrict number of uses of the account
- Literature: PAKE: Password Authenticated Key Establishment

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Entity authentication in practice

- Phishing – mutual authentication
- Forward credentials - biometry
- Interrupt after initial authentication – authenticated key establishment
- Mafia fraud – distance bounding
- Protocol errors – check that local device authentication is linked to entity authentication protocol (example: EMV)

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Mutual authentication


- Phishing is impersonating of the verifier (e.g. the bank)
- Most applications need entity authentication in two directions
- !! This is not complete the same as 2 parallel unilateral protocols for entity authentication

2 stage authentication

- Local: user to device
- Device to rest of the world

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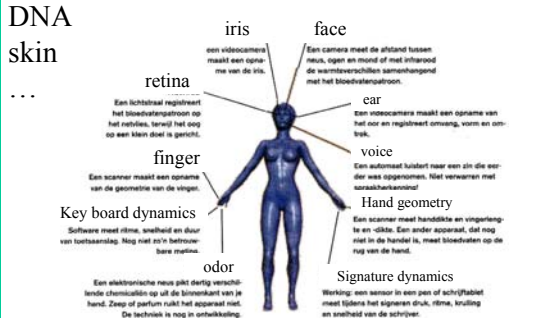
Biometry



- Based on our unique features
- Identification or verification
 - Is this Alice?
 - Check against watchlist
 - Has this person ever registered in the system?

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Some unique features



DNA
skin
retina
iris
face
ear
voice
Hand geometry
Signature dynamics
odor
finger
Key board dynamics

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Biometric procedures

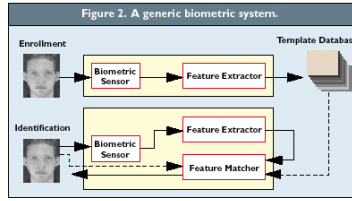


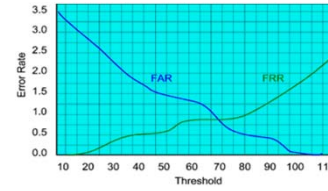
Figure 2. A generic biometric system.

- Registration
- Template extraction
- Measurement
- Processing
- Template matching
- Link with applications

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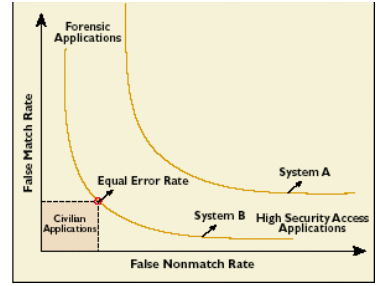
Robustness/performance

- Performance evaluation
 - False Acceptance Ratio or False Match Rate
 - False Rejection Ratio or False Non-Match Rate
- Application dependent



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Robustness/performance (2)

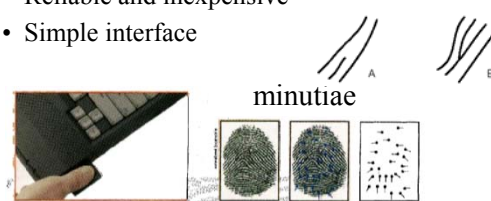


Forensic Applications
 Civilian Applications
 High Security Access Applications
 System A
 System B
 Equal Error Rate

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Fingerprint

- Used for PC/laptop access
- Widely available
- Reliable and inexpensive
- Simple interface



minutiae

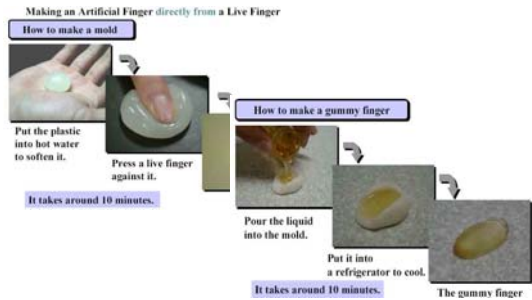
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Fingerprint (2)

- Small sensor
- Small template (100 bytes)
- Commercially available
 - Optical/thermal/capacitive
 - Liveness detection
- Problems for some ethnic groups and some professions
- Connotation with crime

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Fingerprint (3): gummy fingers



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Hand geometry

- Flexible performance tuning
- Mostly 3D geometry
- Example: 1996 Olympics



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Voice recognition

- Speech processing technology well developed
- Can be used at a distance
- Can use microphone of our gsm
- But tools to spoof exist as well
- Typical applications: complement PIN for mobile or domotica

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Iris Scan

- No contact and fast
- Conventional CCD camera
- 200 parameters
- Template: 512 bytes
- All ethnic groups
- Reveals health status



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Retina scan

- Stable and unique pattern of blood vessels
- Invasive
- High security



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Manual signature

- Measure distance, speed, accelerations, pressure
- Familiar
- Easy to use
- Template needs continuous update
- Technology not fully mature



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Facial recognition

- User friendly
- No cooperation needed
- Reliability limited
- Robustness issues
 - Lighting conditions
 - Glasses/hair/beard/...



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Comparison

Feature	Uniqueness	Permanent	Performance	Acceptability	Spoofing
Facial	Low	Average	Low	High	Low
Fingerprint	High	High	High??	Average	High??
Hand geometry	Average	Average	Average	Average	Average
Iris	High	High	High	Low	High
Retina	High	Average	High	Low	High
Signature	Low	Low	Low	High	Low
Voice	Low	Low	Low	High	Low

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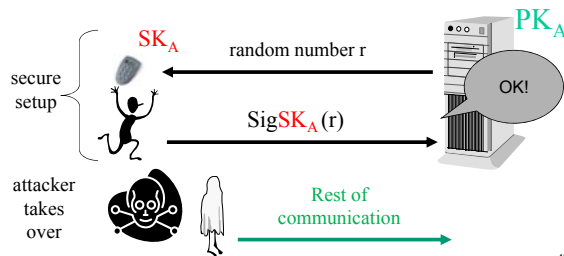
Biometry: pros and cons

- Real person
- User friendly
- Cannot be forwarded
- Little effort for user
- Privacy (medical)
- Intrusive?
- Liveliness?
- Cannot be replaced
- Risk for physical attacks
- Hygiene
- Does not work everyone, e.g., people with disabilities
- Reliability
- No cryptographic key

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Keeping authenticity alive

- Establish who someone is
- Establish that this person is active/liveliness
- But what if the connection is broken after the initial phase?

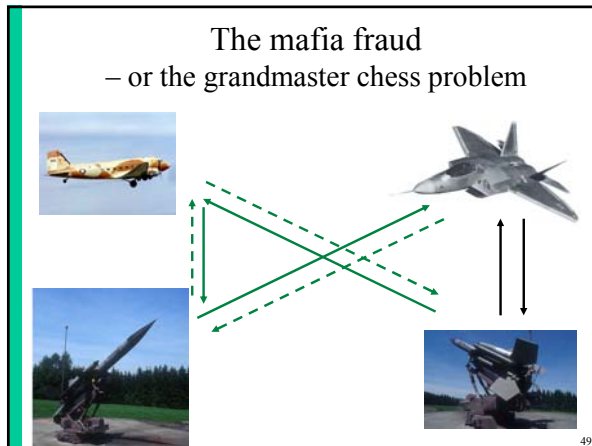


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Solution

- Authenticated key agreement
- Run a mutual entity authentication protocol
- Establish a key
- Encrypt and authenticate all information exchanged using this key

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Location-based authentication

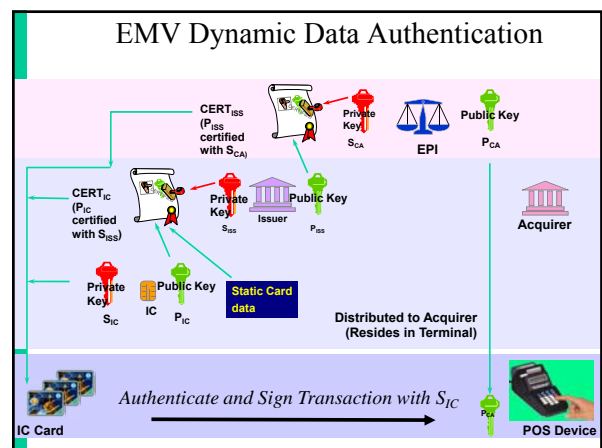
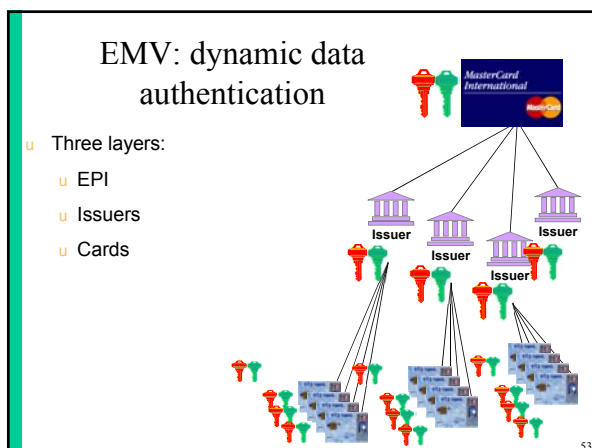
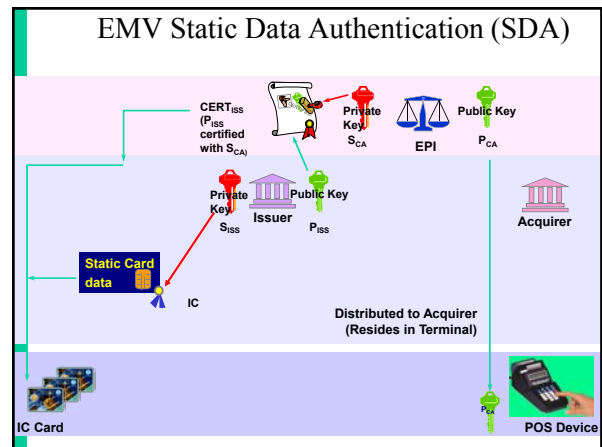
- Distance bounding: try to prove that you are physically close to the verifier
- Other uses of “location”
 - Dial-back: can be defeated using fake dial tone
 - IP addresses and MAC addresses can be spoofed
 - Mobile/wireless communications: operator knows access point, but how to convince others?
 - Trusted GPS: Galileo?

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Authentication with device

- E.g. smart card, secure login token
- Needs 2 stages
 - Local: user to device
 - Device to rest of the world
- Are these 2 stages connected properly?

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Warning about EMV

<http://www.cl.cam.ac.uk/research/security/banking/nopin/oakland10chipbroken.pdf>

- **EMV PIN verification “wedge” vulnerability** S.J. Murdoch, S. Drimer, R. Anderson, M. Bond, IEEE Security & Privacy 2010

Normal PIN check

Fraudulent PIN check

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Guidelines

NIST Special Publication 800-63 Version 1.0.2 (2006):
Electronic Authentication Guideline: identifies four
levels of assurance

http://csrc.nist.gov/publications/nistpubs/800-63/SP800-63V1_0_2.pdf

See <http://csrc.nist.gov/publications/PubsSPs.html>
for about 120 Special Publications (800 Series) from NIST on
computer security and cryptography

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