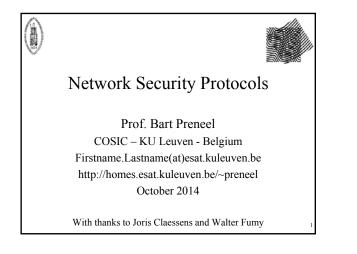
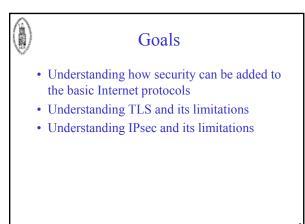
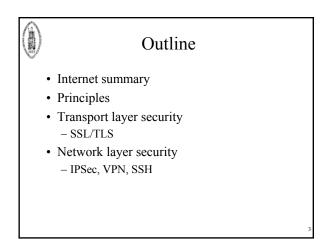
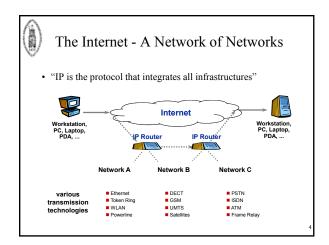
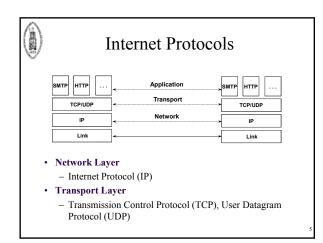
October 2014

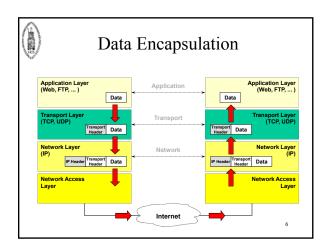


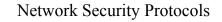


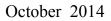


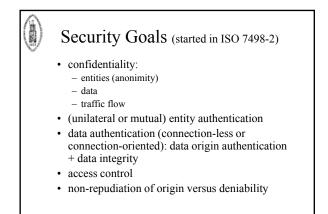


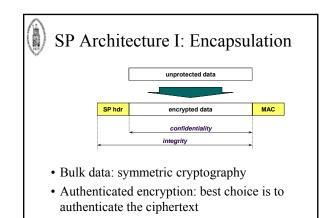


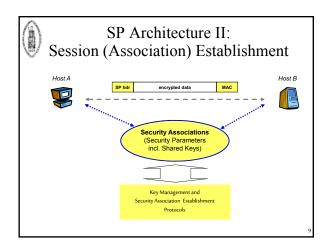


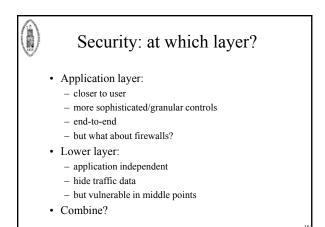


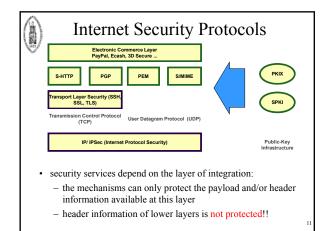


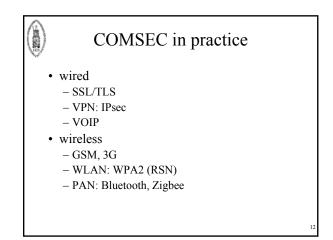




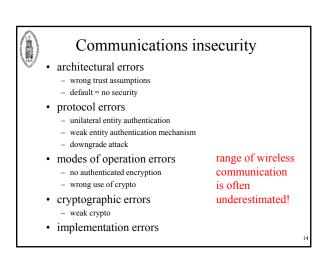




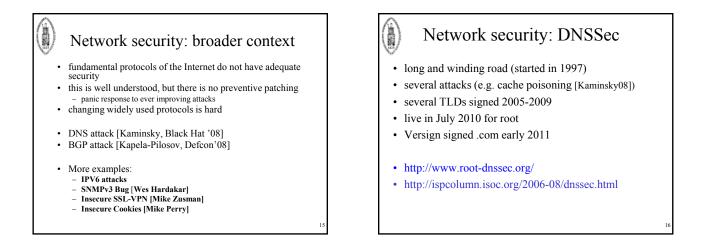




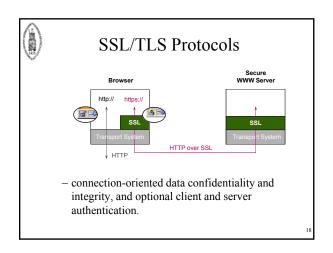
A DESCRIPTION OF A DESC	COMSEC				
25 10		Confidentiality	Data authentication	Entity authentication	
	1 G (analog)				Not
	2 G (GSM)	weak		unilateral	end
	3G				b to end
	WLAN				
	TLS			unilateral	
	IPsec		optional 🛞		
	Skype	not open	not open	not open/meet in the middle attack	13

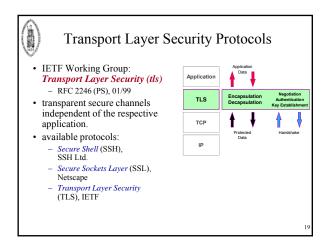


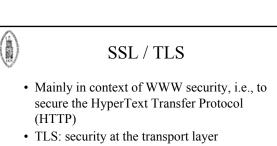
October 2014



Transport layer security SSL / TLS

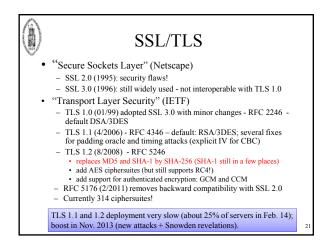


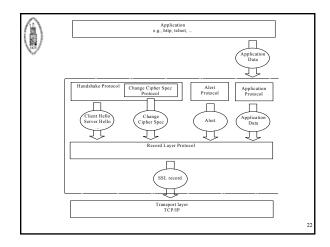


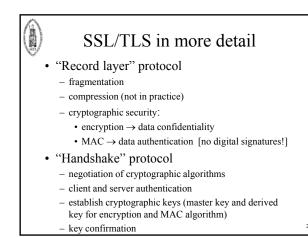


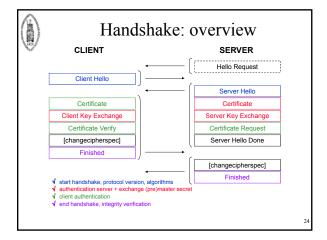
October 2014

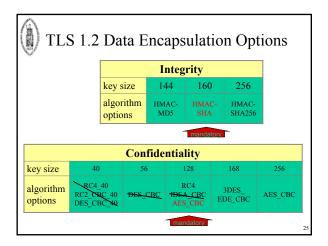
- can be used (and is intended) for other applications too (IMAP, telnet, ftp, ...)
- end-to-end secure channel, but nothing more...
- data is only protected during communication
- no non-repudiation!



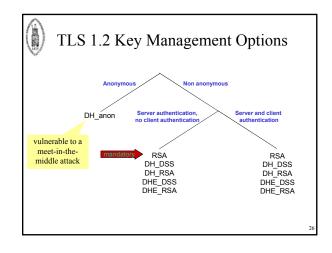


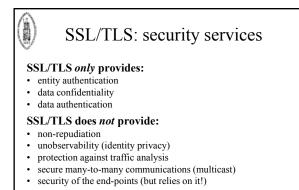






October 2014

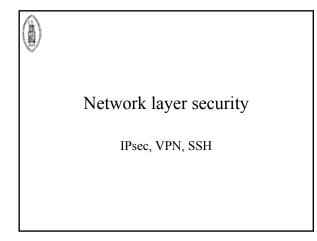


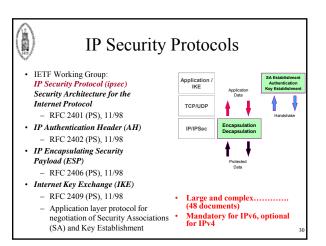


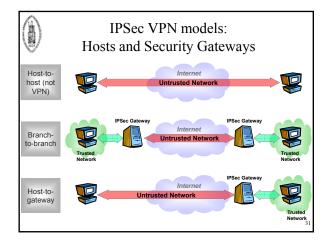
TLS in the future

27

- Reduce the number of cipher suites
- Authenticated encryption gains popularity: - AES-GCM
 - Chacha20 with Poly1305AES
- TLS 2.0: mandatory encryption for httpv2.0?
- Identity protection (cf. IPsec)
- Backward compatibility remains very important because of huge installed base





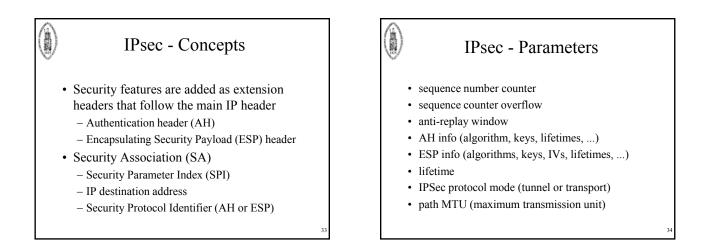


IPsec - Security services Access control Connectionless integrity

October 2014

- Data origin authentication
- Rejection of replayed packets (a form of partial sequence integrity)
- Confidentiality

• Limited traffic flow confidentiality



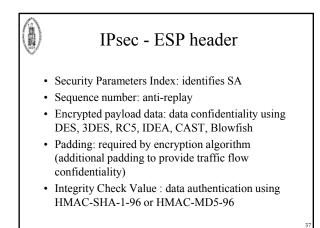
IKE Algorithm Selection Mandatory Algorithms

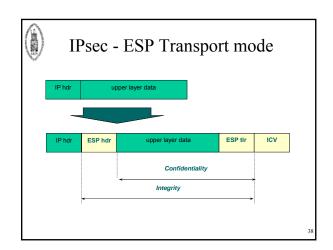
Algorithm Type	IKE v1	IKE v2	
Payload Encryption	DES-CBC	AES-128-CBC	
Payload Integrity	HMAC-MD5 HMAC-SHA1	HMAC-SHA1	
DH Group	768 Bit	1536 Bit	
Transfer Type 1 (Encryption)	ENCR_DES_CBC	ENCR_AES_128_CBC	
Transfer Type 2 (PRF)	PRF_HMAC_SHA1 [RFC2104]	PRF_HMAC_SHA1 [RFC2104]	
Transfer Type 3 (Integrity)	AUTH_HMAC_SHA1_96 [RFC2404]	AUTH_HMAC_SHA1_96 [RFC2404]	

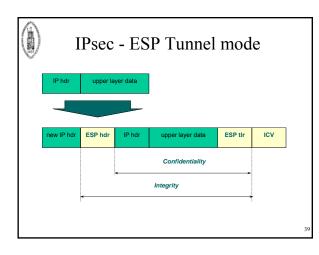
IPsec - Modes

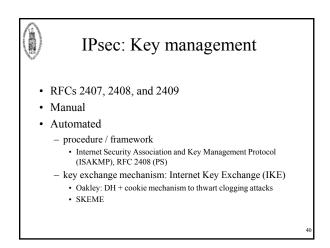
- Transport (*host-to-host*)
 - ESP: encrypts and optionally authenticates IP payload, but not IP header
 - AH: authenticates IP payload and selected portions of IP header
- Tunnel (between security gateways)
 - after AH or ESP fields are added, the entire packet is treated as payload of new outer IP packet with new outer header
 - used for VPN

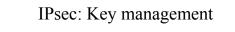
October 2014











- IKE defines 5 exchanges
 - Phase 1: establish a secure channel
 - Main mode

- Aggressive mode
- Phase 2: negotiate IPSEC security association
- Quick mode (only hashes, PRFs)
- Informational exchanges: status, new DH group
- based on 5 generic exchanges defined in ISAKMP
- cookies for anti-clogging

IPsec: Key management

- protection suite (negotiated)
 - encryption algorithm
 - hash algorithm
 - authentication method:
 - preshared keys, DSA, RSA, encrypted nonces
 - Diffie Hellman group: 5 possibilities

IKE v2 - RFC Dec 2005

- IKEv1 implementations incorporate additional functionality including features for NAT traversal, legacy authentication, and remote address acquisition, not documented in the base documents
- Goals of the IKEv2 specification include

問題

戲習

識習

- to specify all that functionality in a single document
- to simplify and improve the protocol, and to fix various problems in IKEv1 that had been found through deployment or analysis
- IKEv2 preserves most of the IKEv1 features while redesigning the protocol for efficiency, security, robustness, and flexibility

October 2014

IKE v2 Initial Handshake (1/2)

- Alice and Bob negotiate cryptographic algorithms, mutually authenticate, and establish a session key, creating an IKE-SA
- Usually consists of two request/response pairs
 - The first pair negotiates cryptographic algorithms and does a Diffie-Hellman exchange
 - The second pair is encrypted and integrity protected with keys based on the Diffie-Hellman exchange

IKE v2 Initial Handshake (2/2)

- · Second exchange
 - divulge identities
 - prove identities using an integrity check based on the secret associated with their identity (private key or shared secret key) and the contents of the first pair of messages in the exchange
 - establish a first IPsec SA ("child-SA") is during the initial IKE-SA creation

間間

IPsec Overview

- · much better than previous alternatives
- · IPsec documents hard to read
- committee design: too complex
 ESP in Tunnel mode with authenticated encryption probably sufficient
- simplify key management
- clarify cryptographic requirements
- ...and thus difficult to implement (securely)
- avoid encryption without data authentication

Concluding comments

- IPsec is really transparent, SSL/TLS only conceptually, but not really in practice
- SSH, PGP: stand-alone applications, immediately and easy to deploy and use
- Network security: solved in principle but – many implementation issues
 - complexity creates security weaknesses
- Application and end point security: more is needed!

More information (1)

- William Stallings, *Cryptography and Network Security - Principles and Practice*, Fifth Edition, 2010
- N. Doraswamy, D. Harkins, *IPSec (2nd Edition)*, Prentice Hall, 2003 (outdated)
- Erik Rescorla, SSL and TLS: *Designing and Building Secure Systems*, Addison-Wesley, 2000.
- IETF web site: www.ietf.org
 e.g., IETF-TLS Working Group http://www.ietf.org/html.charters/tls-charter.html

More information (2)

- Jon C. Snader, VPNs Illustrated: Tunnels, VPNs, and IPsec, Addison-Wesley, 2005
- Sheila Frankel, *Demystifying the IPsec Puzzle*, Artech House Computer Security Series, 2001
- Anup Gosh, E-Commerce Security, Weak Links, Best Defenses, Wiley, 1998
- Rolf Oppliger, Security Technologies for the World Wide Web, Artech House Computer Security Series 1999
- W3C Security (incl WWW Security FAQ) http://www.w3.org/Security/